

# Spacetime Programming

Synchron 2016

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# Constraint programming

## Holy grail of computing

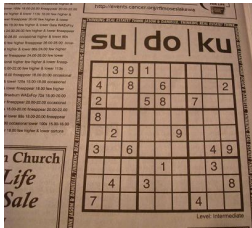
- ▶ Declarative paradigm for solving combinatorial problems.
- ▶ We state the problem and let the system solve it for us.



# Successful paradigm

## Applications

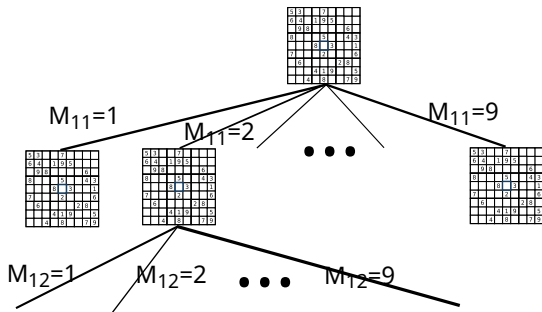
It has a lot of different applications ranging from Sudoku solving, scheduling, packing, musical orchestration...



# How to find a solution?

## NP-complete nature

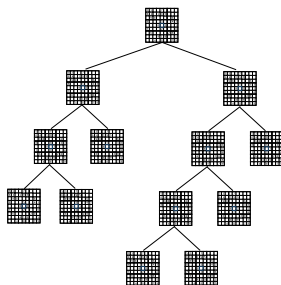
- ▶ Try every combination until we find a solution.
- ▶ The possible combinations are represented in a tree.



# Problem

## Holy grail?

- ▶ Search tree is often too huge to find a solution in a reasonable time.
- ▶ **Search strategies are crucial** for describing how to create and prune the tree and improving efficiency.
- ▶ Search strategies are often problem-dependent so we need to try and test (empirical evaluation).



# State-of-the-art

1. **Languages** (Prolog, MiniZinc,...): Clear and compact description but limited amount of pre-defined strategies.
  2. **Libraries** (Choco, GeCode,...): Highly customizable and efficient but complex software, hard to understand and time-consuming.
- ▶ Composing strategies is impossible or hard in both cases.

**Lack of abstraction** for expressing, composing and extending search strategies.

## Synchronous languages provide the needed abstraction!

- ▶ We propose *spacetime programming*, a language abstraction for expressing search strategies.
- ▶ Based on Esterel (without the reaction to absence).
- ▶ **Execution**: One node of the tree processed per instant.
- ▶ **Nondeterministic** operator for specifying the branches of the tree.



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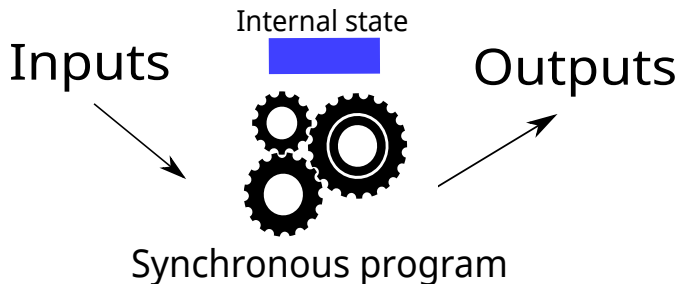
# Spacetime programming

Spacetime programming = Synchronous programming + Search strategy.

- ▶ Search strategies as synchronous processes.
- ▶ **Composition** of strategies with the parallel operator.
  - ▶ `par s1 || s2 end`
  - ▶ Easy experiment: plugging in and out strategies.
- ▶ **Communication** between strategies in the deterministic framework of the synchronous paradigm.

# Synchronous programming

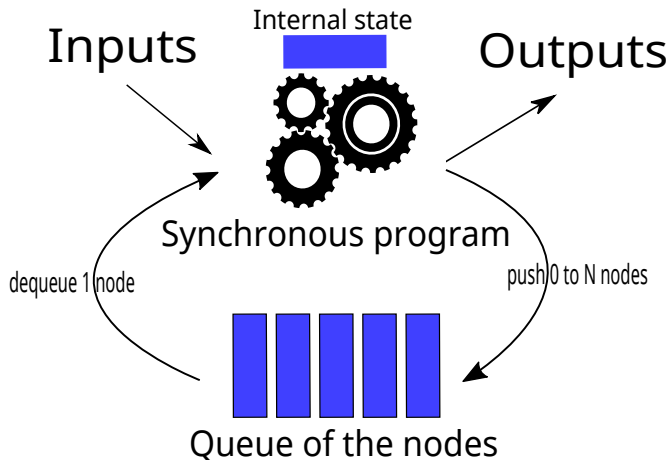
- ▶ In one instant, a synchronous program reacts to inputs and emits outputs.
- ▶ It keeps an internal state of variables and program status.



How to link the synchronous model and search tree?

# Spacetime execution scheme

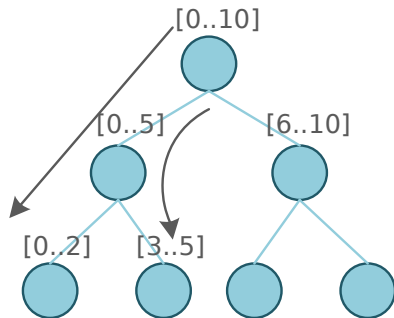
- ▶ The search tree is represented as a queue of nodes.
- ▶ We feed the program with **one node of the tree per instant**.
- ▶ The synchronous program fuels the queue with new nodes.



## Space: Creating the tree

- ▶ `space p || q end` for creating two branches where  $p$  and  $q$  describes children nodes.

```
let x = [0..10];  
loop  
  let mid = middle_value(x);  
  space  
    || x ← [lb(x)..mid-1]  
    || x ← [mid..ub(x)]  
  end  
  pause  
end
```



# Internal state

- ▶ We can use the internal state for maintaining global information to the tree.
- ▶ For example, for maintaining statistics such as the number of nodes explored.

```
count_nodes ≡  
  nodes ← 1;  
  loop  
    pause;  
    nodes ← (pre nodes) + 1;  
  end
```

# Spacetime attribute

## Problem

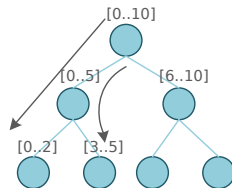
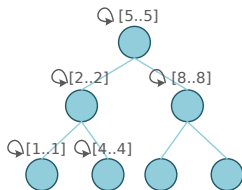
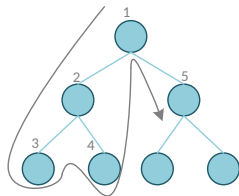
How to differentiate between variables in internal state and onto the queue?

We use a spacetime attribute to situate a variable in space and time.

- ▶ **Global:** Variable in one location, global to the search tree (attribute `single_space`).
- ▶ **Local:** Variable in one time, local to one instant (attribute `single_time`).
- ▶ **Backtrackable:** Variable in the queue of nodes (attribute `world_line`).

# Spacetime attribute

```
let x in world_line = [0..10];  
loop  
  let mid in single_time = middle_value(x);  
  space  
  || x ← [lb(x)..mid-1]  
  || x ← [mid..ub(x)]  
end  
pause  
end
```



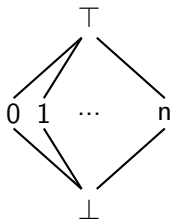


# Variables are complete lattices

- ▶ Every variable is a complete lattice where  $\leftarrow$  is the join operator and  $\text{bot}$  the bottom representing the lack of information.
- ▶ `transient` re-initializes the value to bottom between instants (persistent by default).

```
let transient nodes = bot;
```


```
count_nodes  $\equiv$   
  nodes  $\leftarrow$  1;  
  loop  
    pause;  
    nodes  $\leftarrow$  (pre nodes) + 1;  
  end
```



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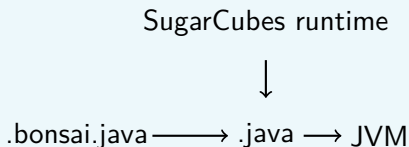
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## Implementation: Bonsai

- ▶ Integration into object-oriented language (Java).
- ▶ Extend the Java syntax with processes and reactive attributes.
- ▶  [github.com/ptal/bonsai](https://github.com/ptal/bonsai)

### Compilation

The compiler acts as a preprocessor from Bonsai to Java.



## SugarCubes (Susini, '01)

SugarCubes is a Java library to program reactive systems with the synchronous paradigm.

- ▶ It provides a set of class combinators for each synchronous instructions.
- ▶ For example, `loop { pause; }` is compiled to `new Loop(new Pause())`.
- ▶ Method `activate()` called at each instant on the combinators.

# Bonsai syntax

- ▶ Must inherits from `Executable` and have a process named `execute` (entry point).
- ▶ Java method call with `~method`.

```
public class ConstraintProblem implements Executable
{
    world_line VarStore domains = bot;
    world_line ConstraintStore constraints = bot;
    proc execute() {
        ~modelChoco(domains, constraints);
        par branching() || propagate() end
    }
    private static void modelChoco(VarStore domains,
        ConstraintStore constraints )
    { ... }
}
```

# Compilation

- ▶ A runtime environment contains all the variables.
- ▶ Programs are created at runtime.

```
public Program execute() {
    return SC.seq(
        new JavaAtom((env) -> {
            VarStore domains = (VarStore) env.var("domains");
            ConstraintStore constraints = (ConstraintStore) env.var(" constraints ");
            modelChoco(domains, constraints);
        }),
        SC.par(
            branching(),
            propagate()
        )
    );
}
```

# Experiments

We validate this approach by replacing the search module of the state-of-the-art constraint solver Choco and comparing the efficiency.

- ▶ We provide a small binding (200 loc) to be able to use Choco inside the language.
- ▶ We implemented the same search strategy in Choco and in Bonsai.
- ▶ Comparison on 3 different constraint problems.

## Experiments

	Choco	SP	$\frac{SP}{Choco}$
	First solution		
Latin square (40)	3.42 s	3.45 s	1
Latin square (50)	8.26 s	9.66 s	1.17
Latin square (60)	19.49 s	23.20 s	1.19
	All solutions		
N-Queens (12)	1.44 s	3.62 s	2.51
N-Queens (13)	6.35 s	16.04 s	2.53
N-Queens (14)	32.10 s	147 s	4.58
	Best solution		
Golomb ruler (9)	0.57 s	1.61 s	2.83
Golomb ruler (10)	1.69 s	6.43 s	3.81
Golomb ruler (11)	24.89 s	135 s	5.42

- ▶ Almost no overhead for finding one solution, factor between 2 and 5 for all and best solution.



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# Conclusion

- ▶ Lack of an abstraction for expressing search strategies.
- ▶ Synchronous language is an ideal abstraction when extended with:
  - ▶ Partial information (lattice-based variable).
  - ▶ Nondeterminism.
- ▶ Working implementation available.
- ▶ Experiments show an acceptable overhead compared to state-of-the-art solvers.

 `github.com/ptal/bonsai`

# Future work

- ▶ **Static analysis** for avoiding the top value.
- ▶ **Interactive constraint system.**
  - ▶ Computer-aided composition with constraints.
- ▶ Queue of nodes directly accessible in the program.
  - ▶ Enables restart-based search strategies such as iterative deepening, limited discrepancy, ...

Thank you for your attention.

